

Representing integer multiplication using BDDs

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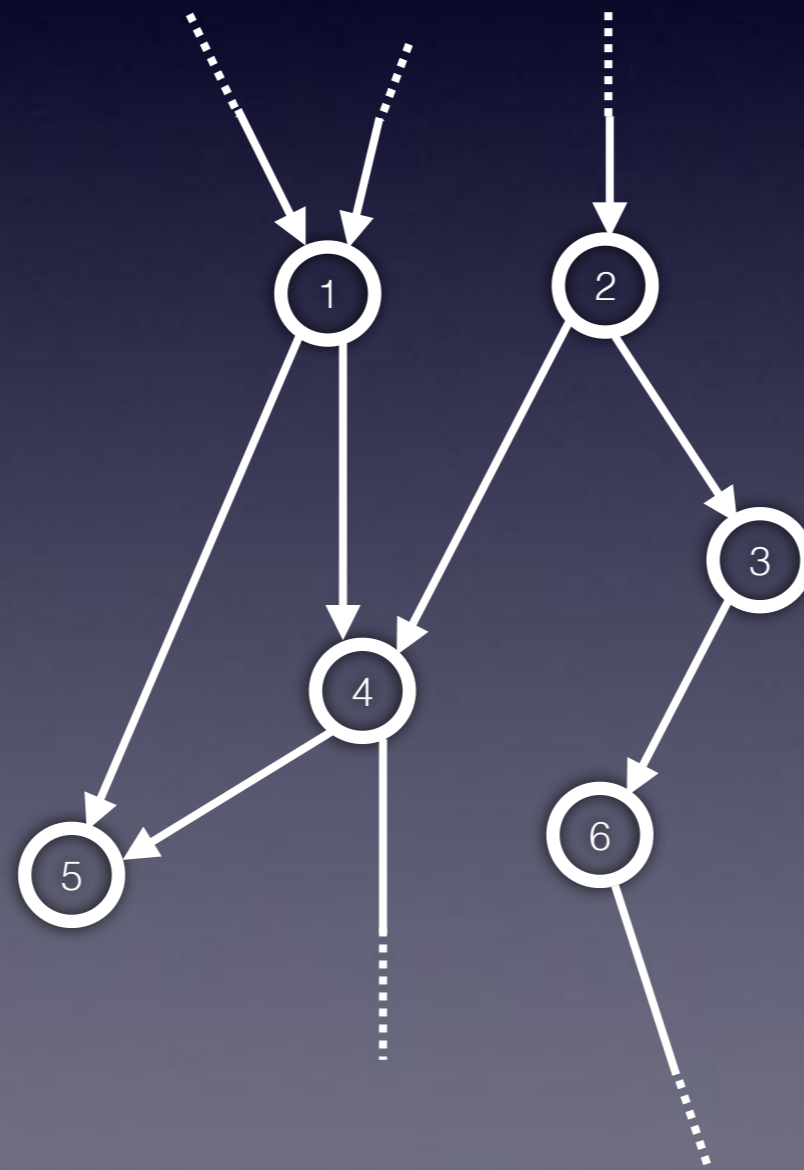
What is a BDD?

Can be defined and understood in different ways

Only one description given here

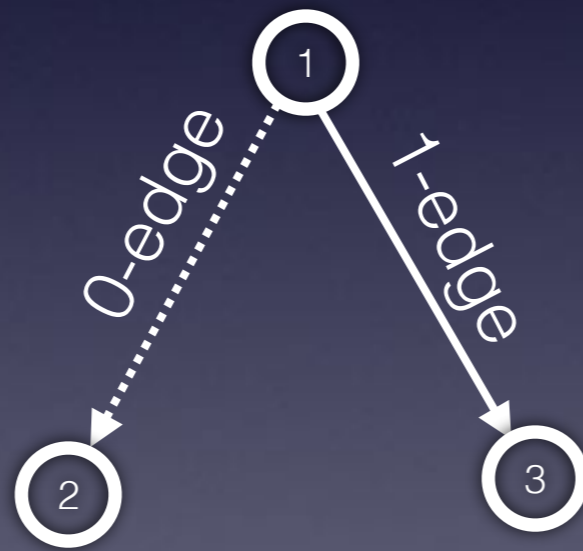
BDD explained

A BDD is a directed acyclic graph, drawn from top to bottom



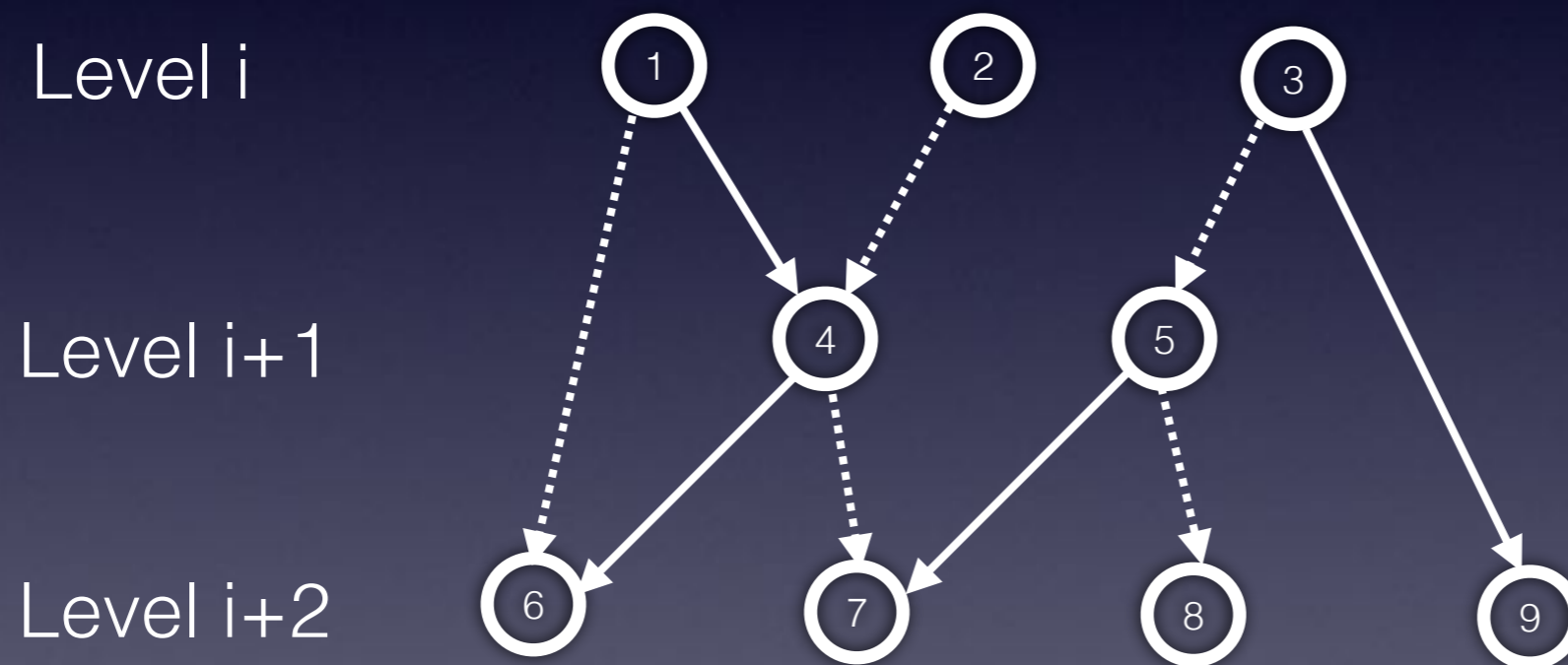
BDD explained

Every node can have two outgoing edges, the 0-edge and the 1-edge



BDD explained

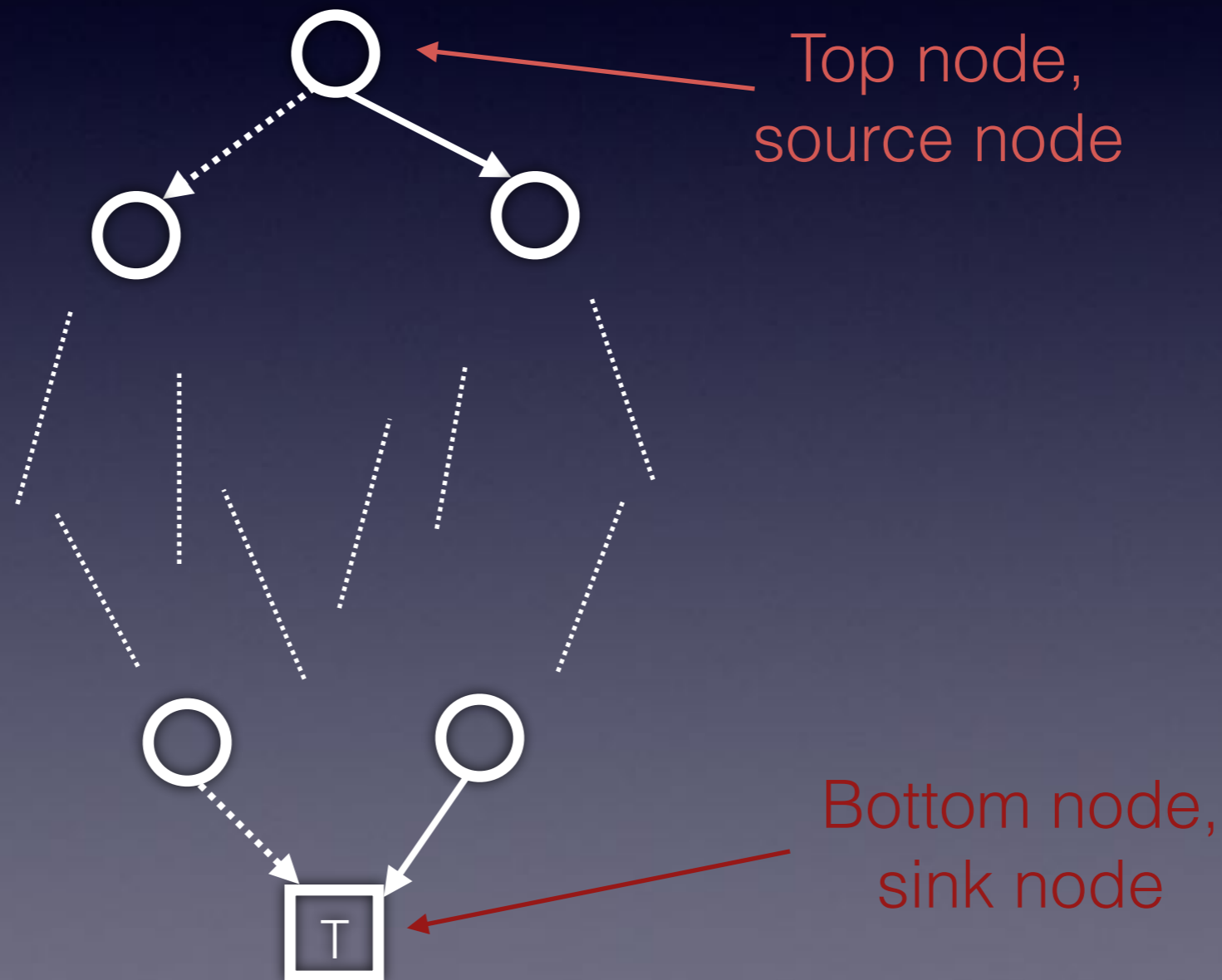
Nodes are arranged in horizontal levels



No edges between nodes on same level

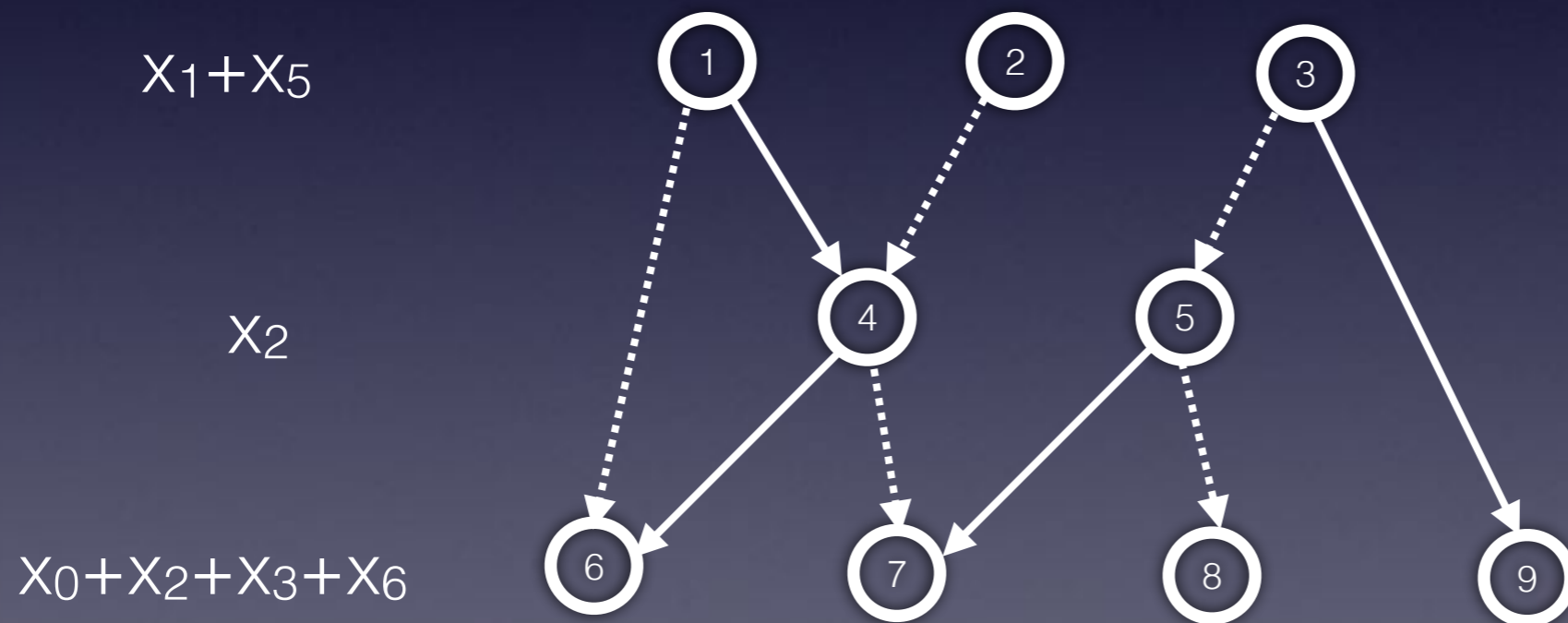
BDD explained

Only one node on highest and lowest levels



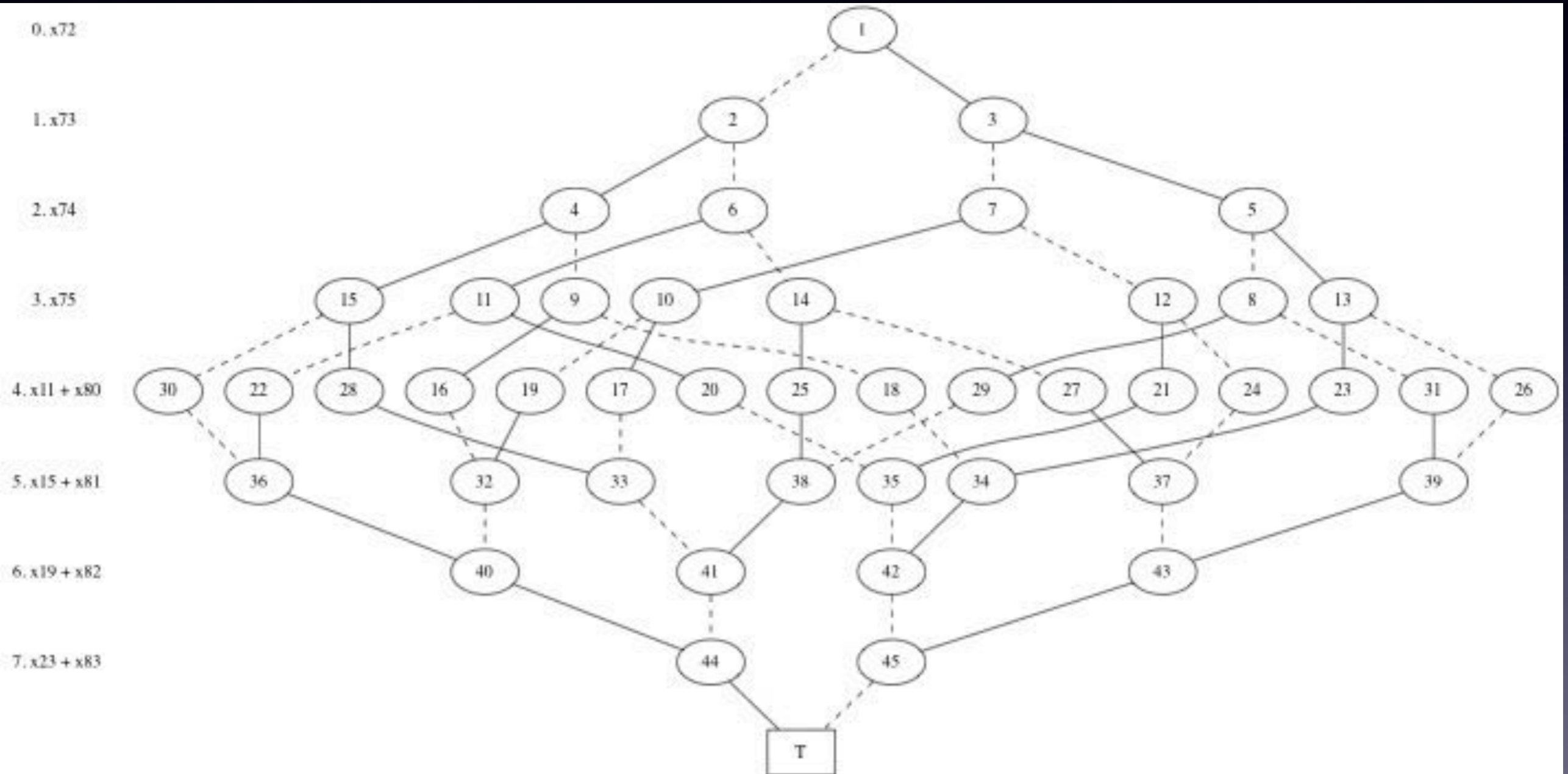
BDD explained

Linear combination of variables associated with each level, except bottom level

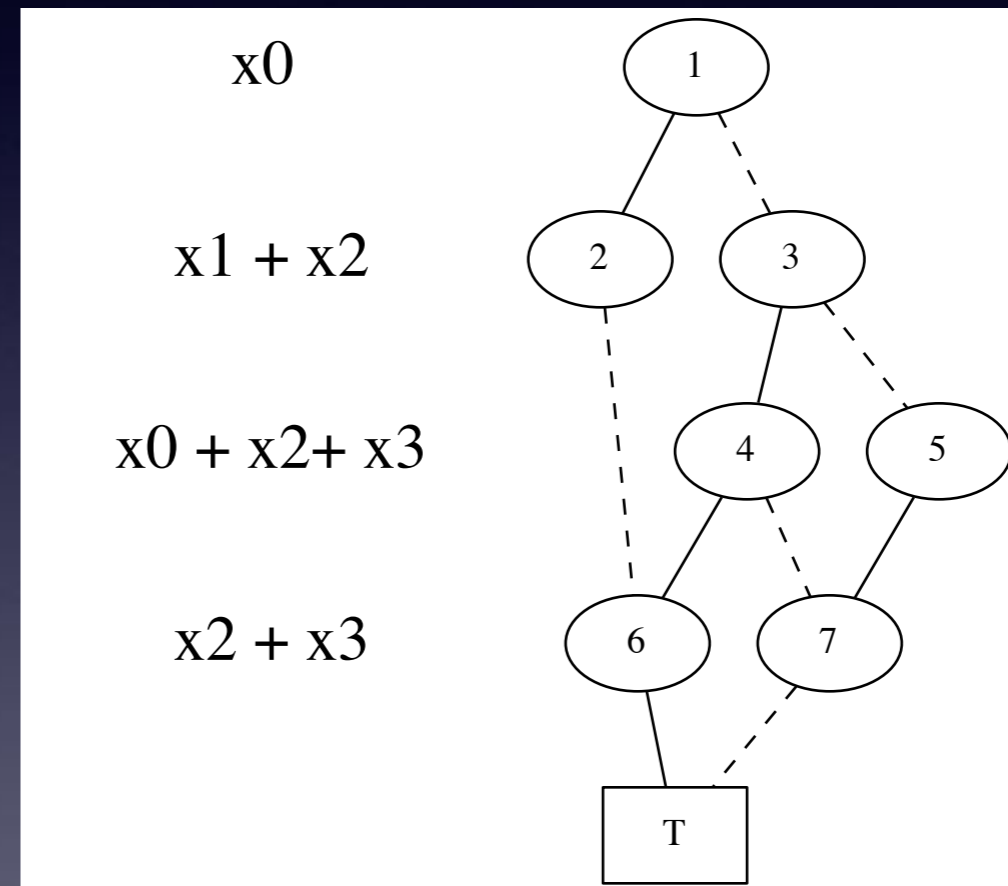
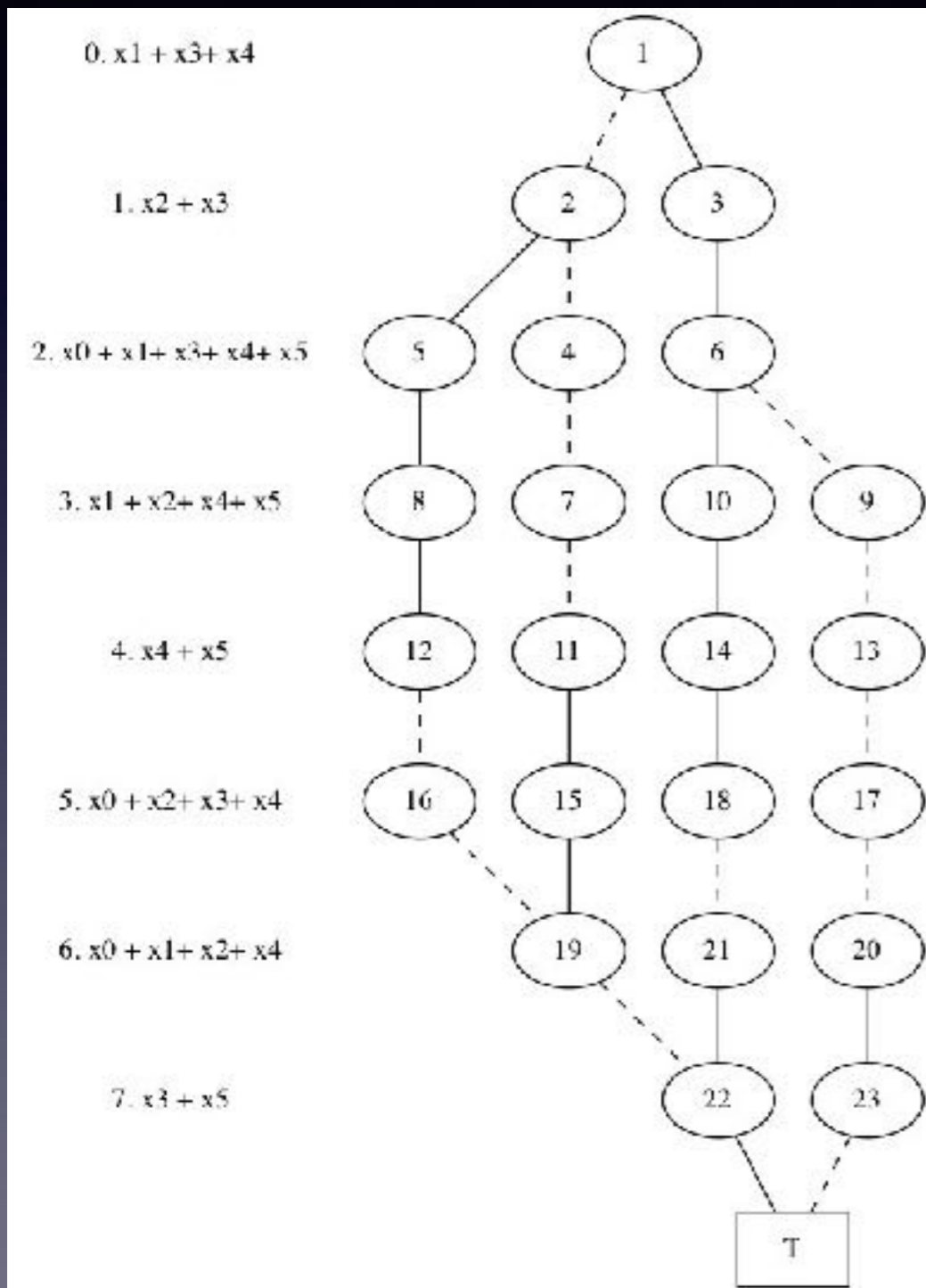


Other descriptions: only single variables associated with levels

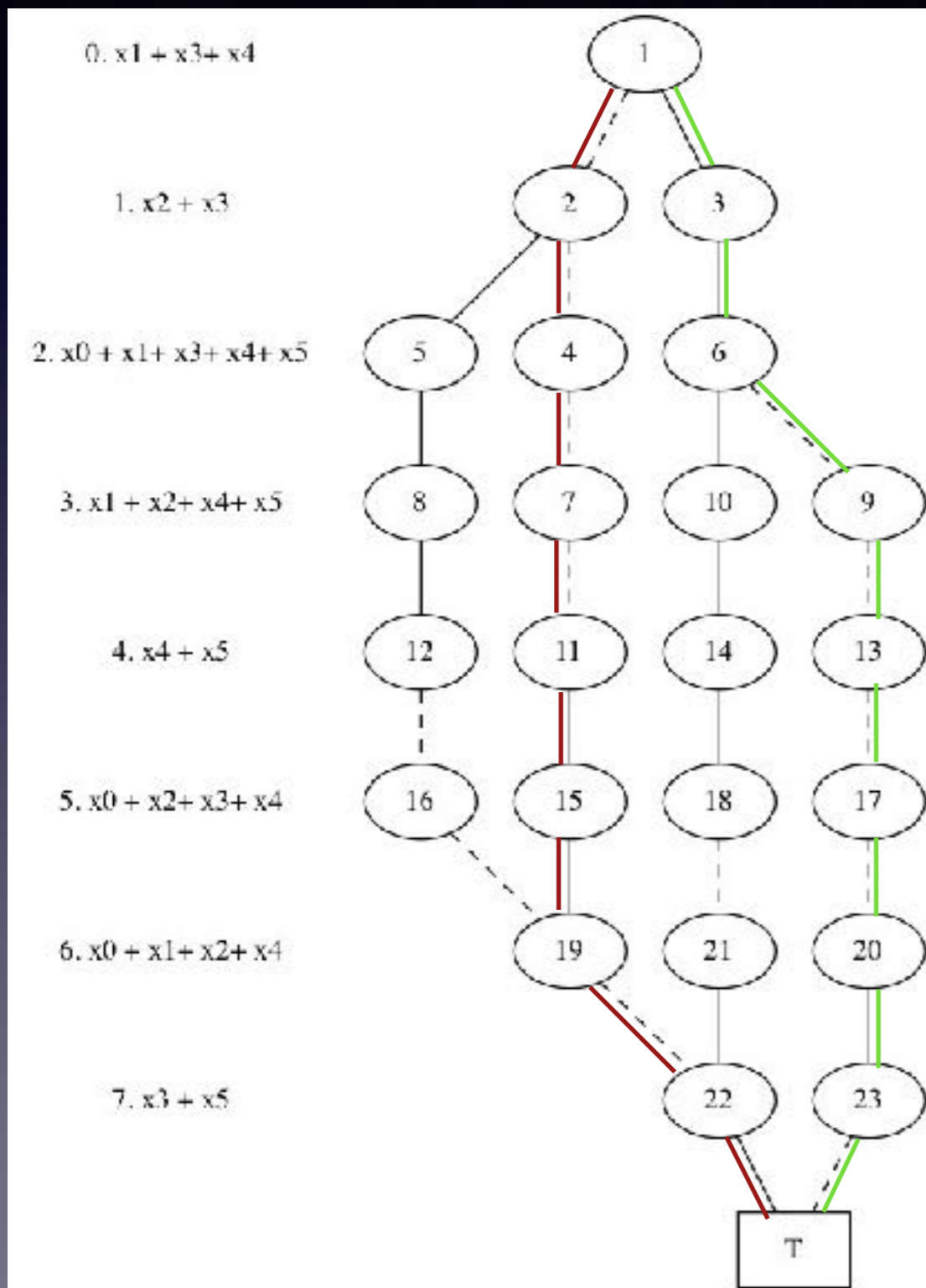
BDD Examples



BDD examples



Paths

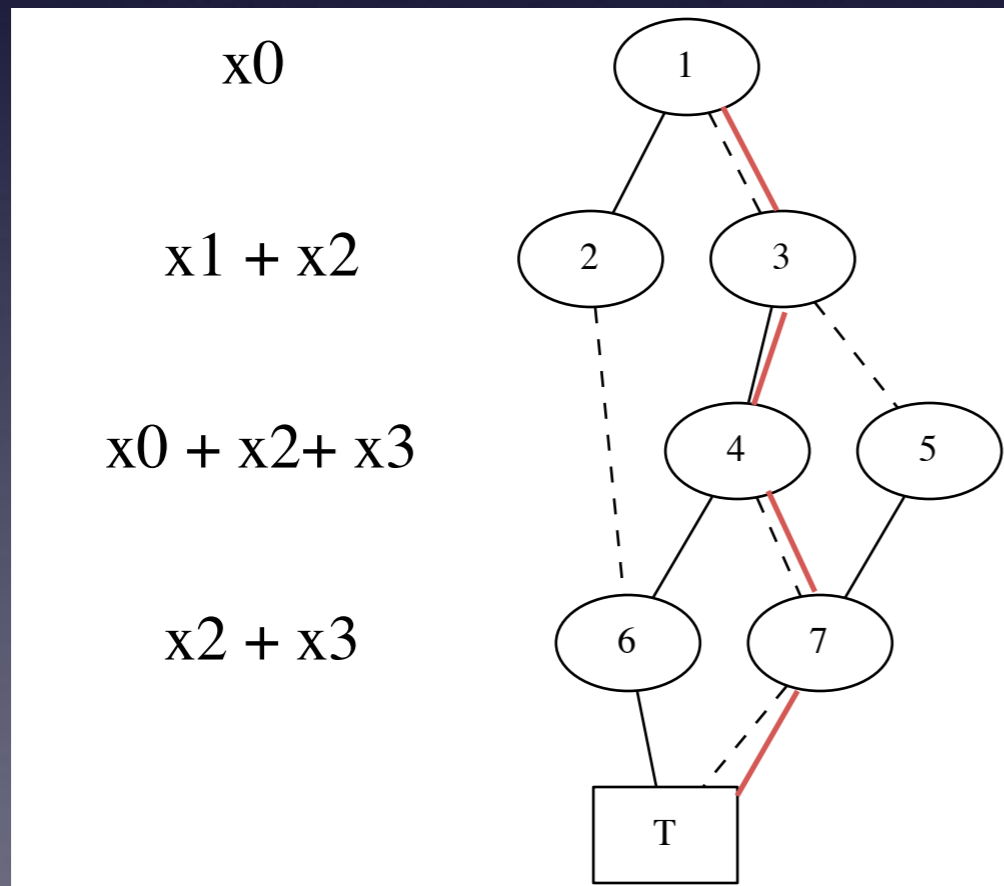


0
0
0
0
0
1
1
0
1

1
1
0
0
0
0
0
1
0

Assigning values

Values in path assigned to linear combinations associated with levels



$$x_0 = 0$$
$$x_1 + x_2 = 1$$
$$x_0 + x_2 + x_3 = 0$$
$$x_2 + x_3 = 0$$

Integer multiplication as BDD

How to construct multiplication BDD?

N, p and q in binary

- RSA modulus $N=pq$
- $p=p_{n-1} p_{n-2} \dots p_2 p_1 p_0$
- $q=q_{n-1} q_{n-2} \dots q_2 q_1 q_0$
- $N=N_{2n-1} N_{2n-2} \dots N_2 N_1 N_0$
- $N_i, p_i, q_i \in \{0, 1\}$

Multiplication

	$(q_{n-1}$	\dots	q_2	q_1	$q_0)$	$(p_{n-1}$	\dots	p_2	p_1	$p_0)$
						p_0q_{n-1}	\dots	p_0q_2	p_0q_1	p_0q_0
+					p_1q_{n-1}	p_1q_{n-2}	\dots	p_1q_1	p_1q_0	
+				p_2q_{n-1}	p_2q_{n-2}	p_2q_{n-3}	\dots	p_2q_0		
\vdots			\ddots	\vdots	\vdots	\vdots	\ddots			
+		$p_{n-1}q_{n-1}$	\dots	$p_{n-1}q_2$	$p_{n-1}q_1$	$p_{n-1}q_0$				
=	N_{2n-1}	N_{2n-2}	\dots	N_{n+1}	N_n	N_{n-1}	\dots	N_2	N_1	N_0
		$2n-2$		$n+1$	n	$n-1$		2	1	0

columns

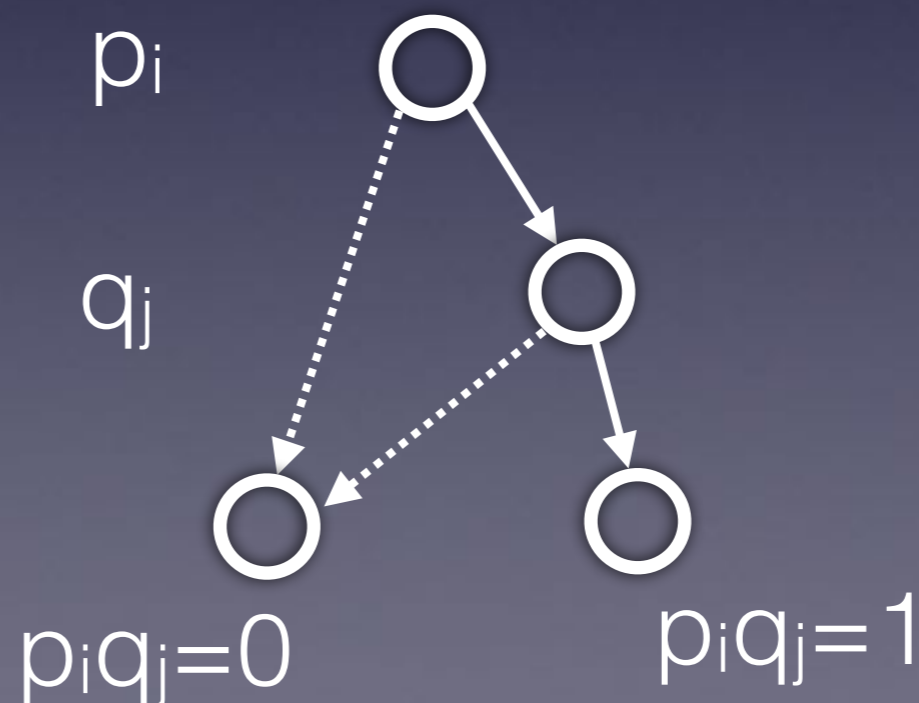
Column k consists of all terms p_iq_j where $i+j=k$

Column k contributes $2^k \sum p_iq_j$ to the value of N

Building multiplication BDD

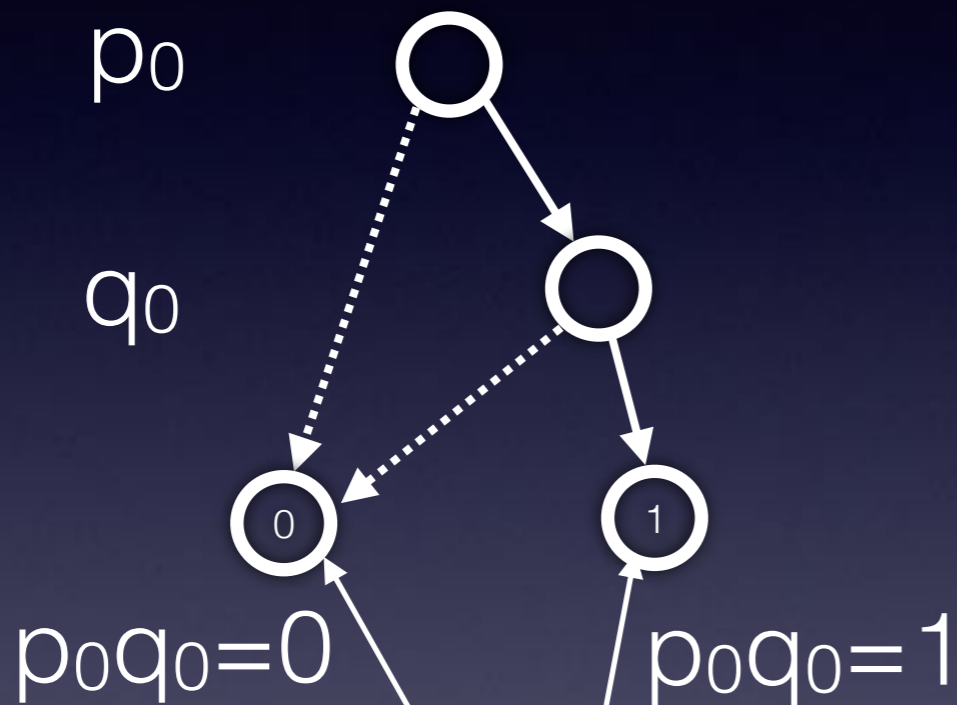
- BDD built column by column, starting with column 0
- One level of nodes constructed from each variable p_i , q_j appearing in column k

Basic building block:



Column 0

$N = \dots 011$



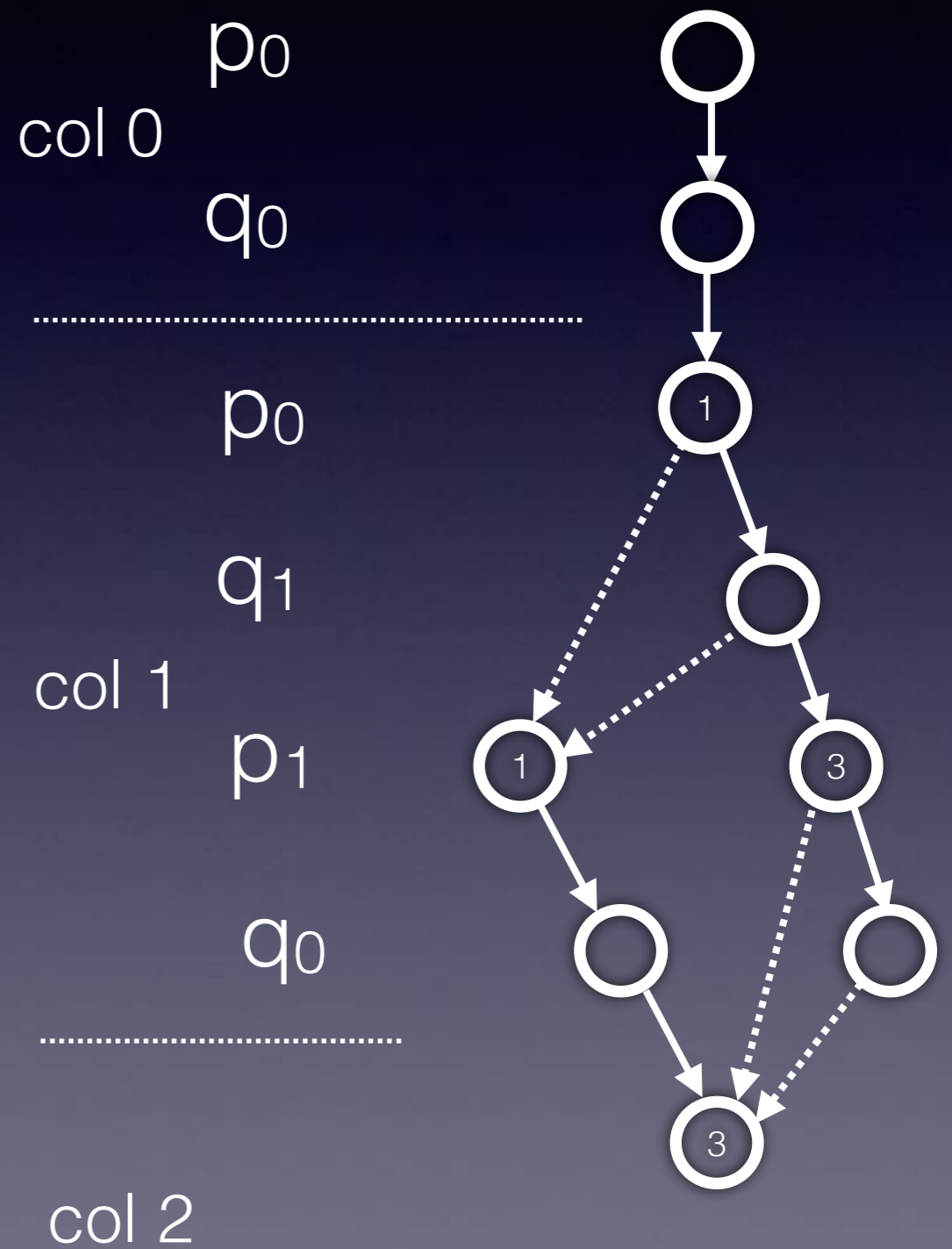
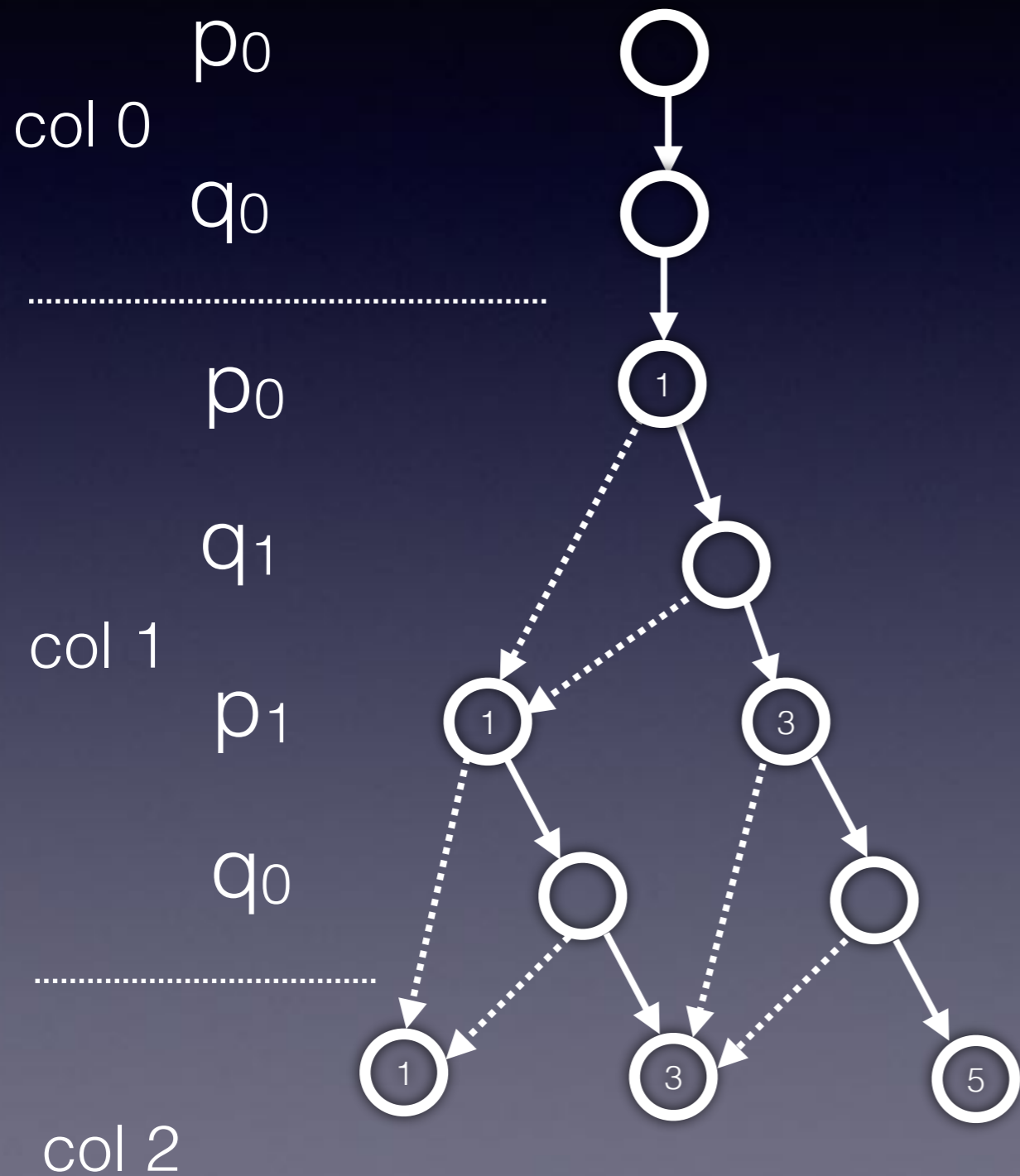
keep track of
value of N
computed so far



When construction of column k
is complete, delete nodes with
values inconsistent with N_k

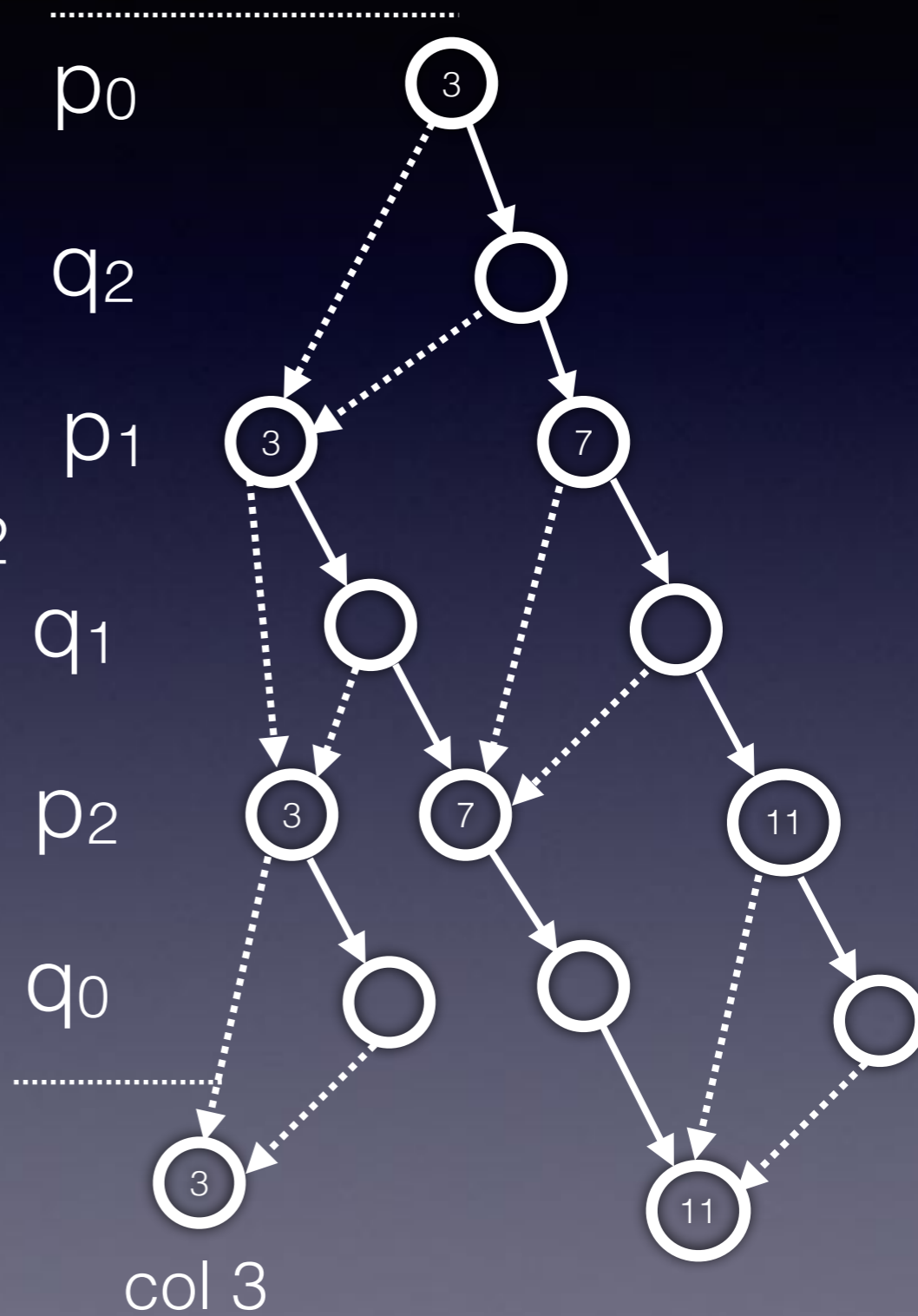
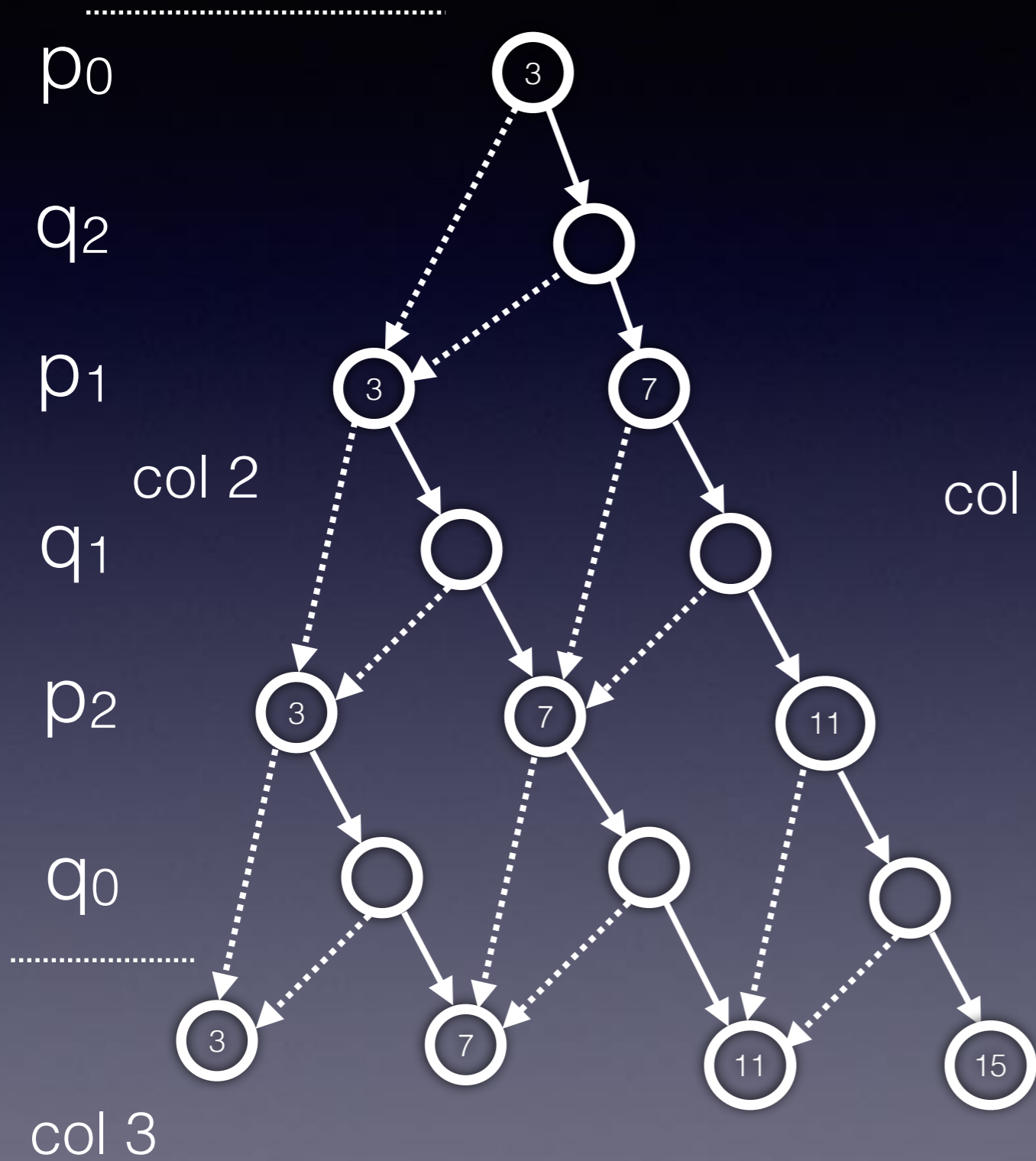
Column 1

$N = \dots 011$

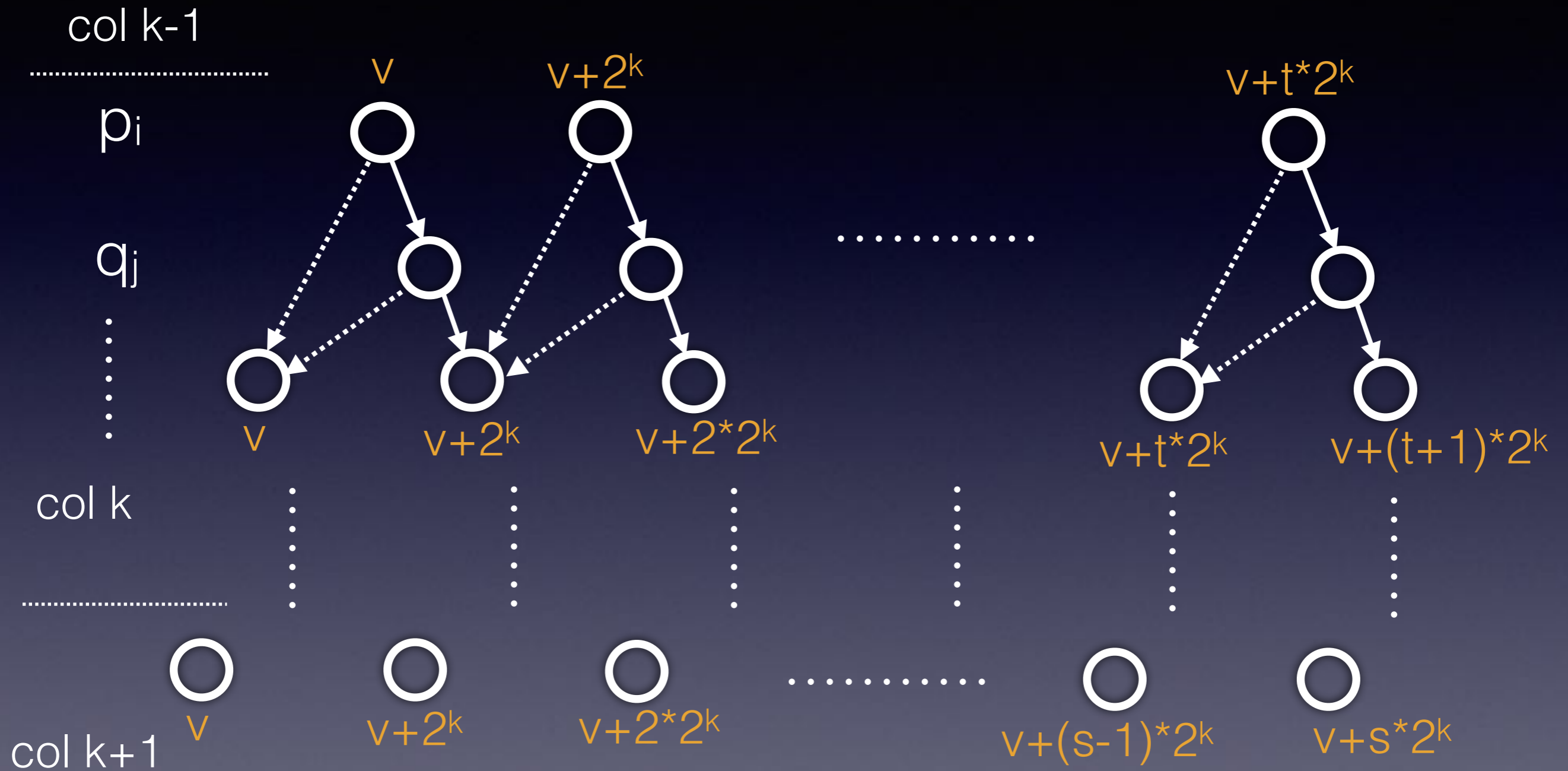


Column 2

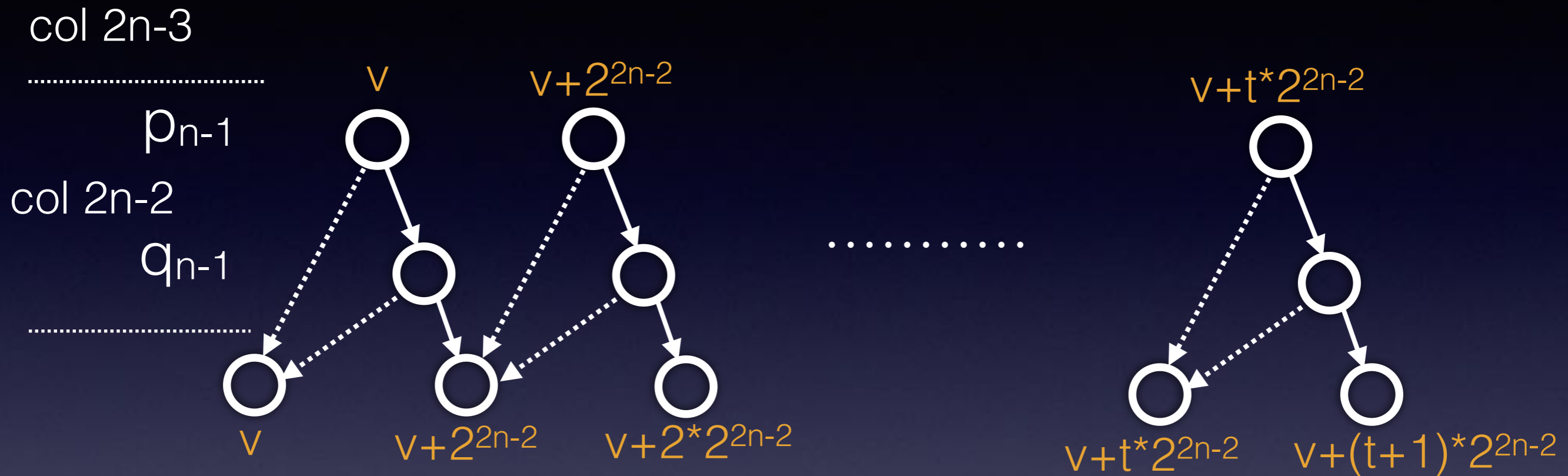
$N = \dots 011$



Column k

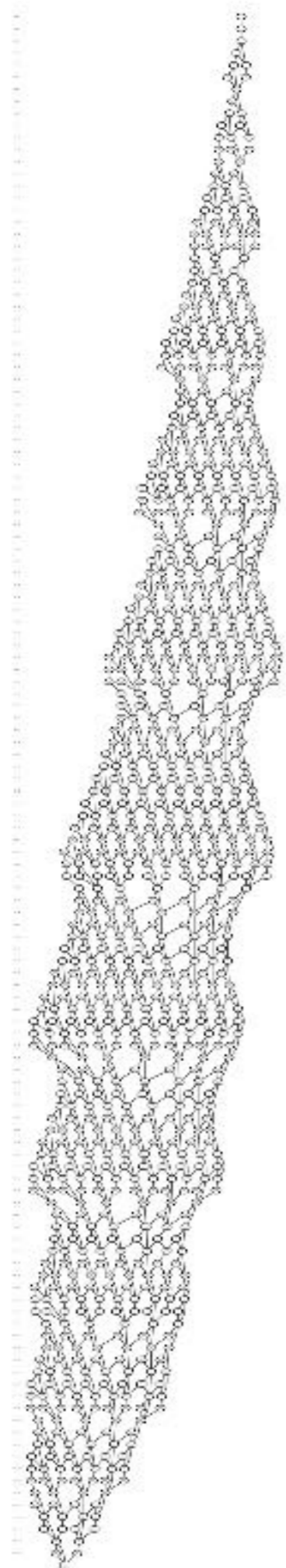


Column $2n-2$



Only keep node with value N .

This becomes the bottom node of the BDD



BDD for multiplication of two 8-bit numbers

$N=31439$

891 nodes

$> 2^{110}$ paths

General properties of multiplication-BDD

- BDD for n-bit numbers p and q has:
 - $2n$ different variables
 - $2n^2$ levels (each variable occurring n times)
 - $\leq 2n^3 - 9/2n + 5$ nodes

Using BDD for factoring

How to find values of p and q ?

Finding lsb's of p and q

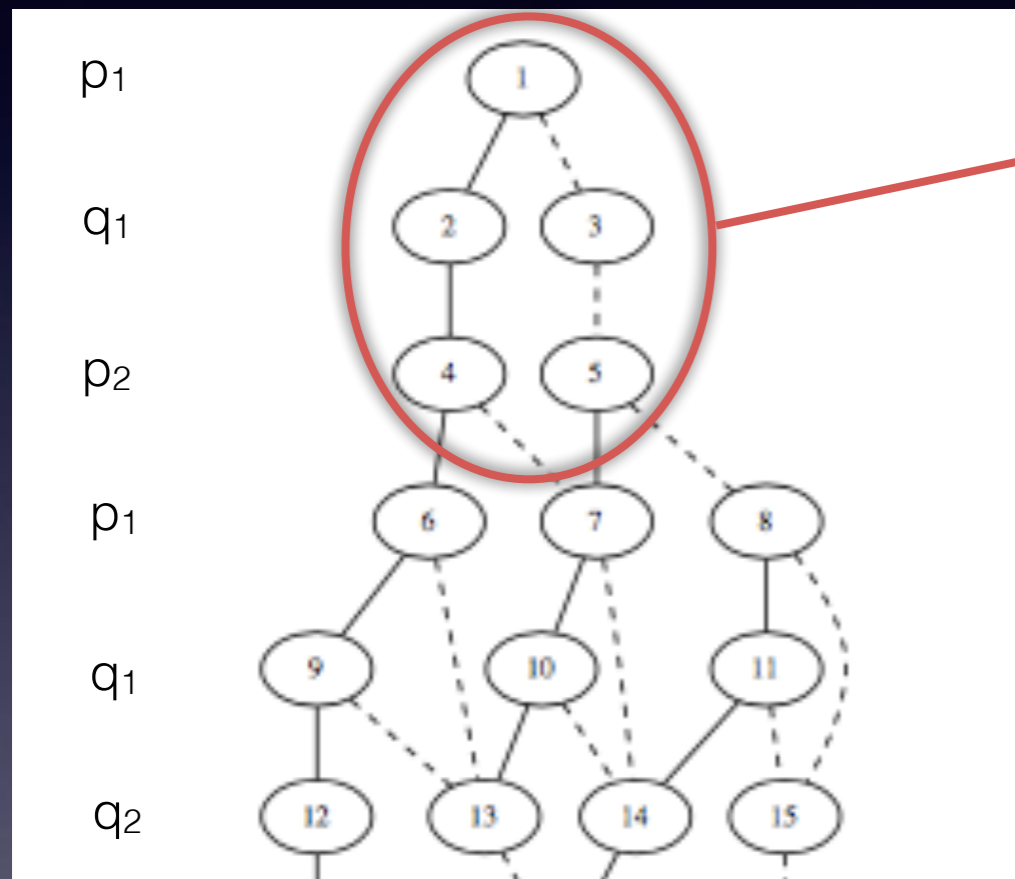
- When N is odd, top of BDD will always be



The two top levels
give $p_0=q_0=1$

Fix $p_0=q_0=1$ in whole BDD by deleting 0-edges
going out from levels with p_0 or q_0

Finding Isb's of p and q



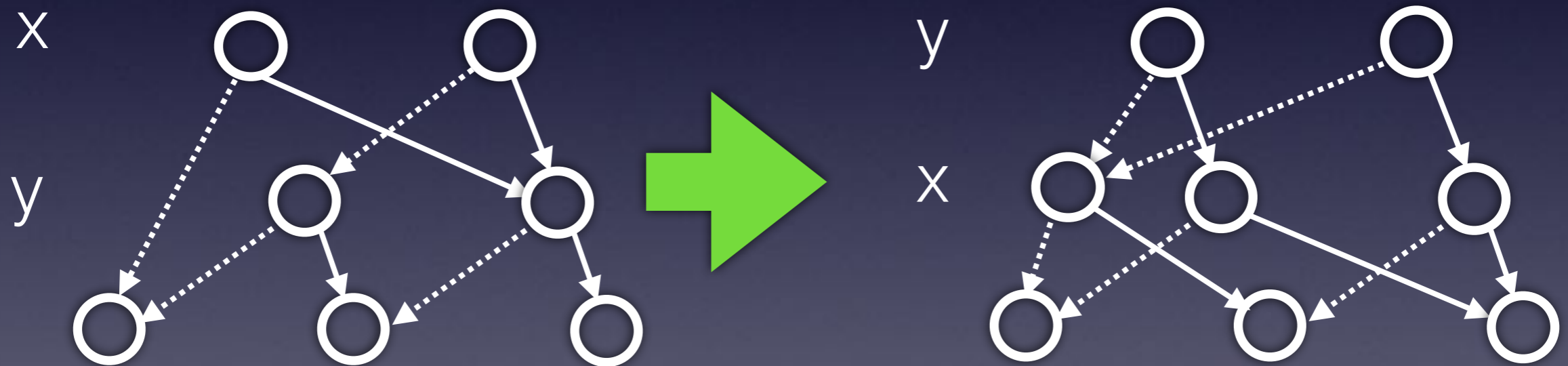
Gives $p_1 = q_1$

Replace q_1 with p_1
on all levels

Depending on value of N , some more linear equations can be extracted from top of BDD

Swapping levels

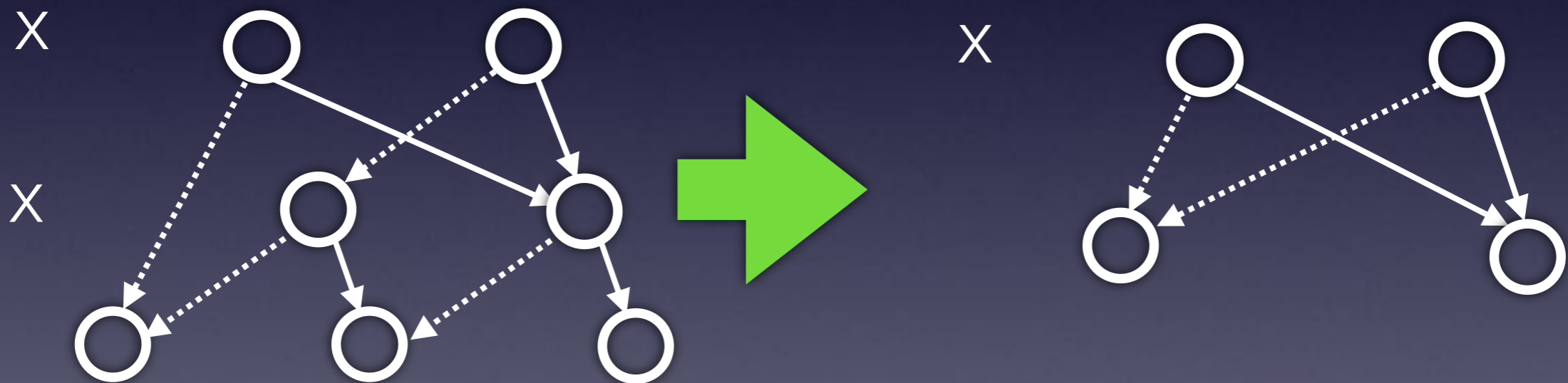
- There exists an algorithm for swapping adjacent levels



- Number of nodes may increase on affected level

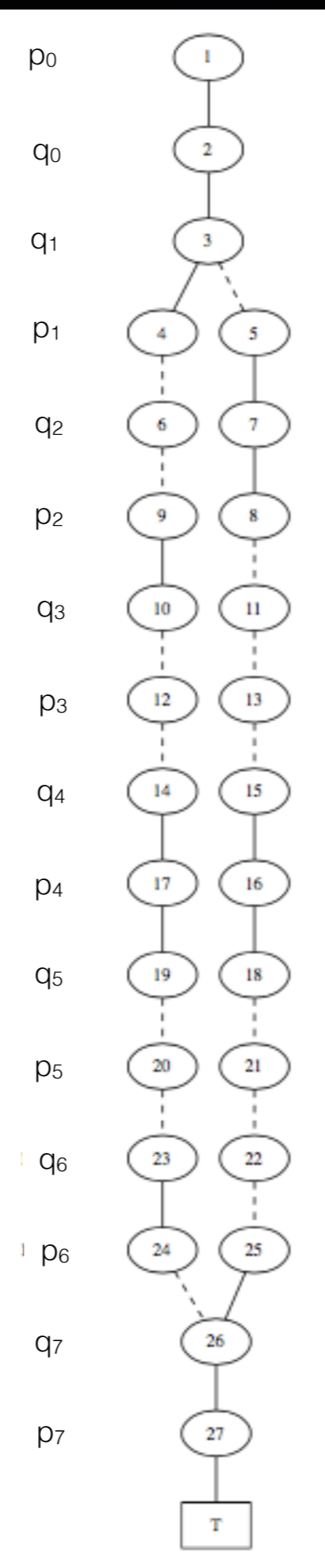
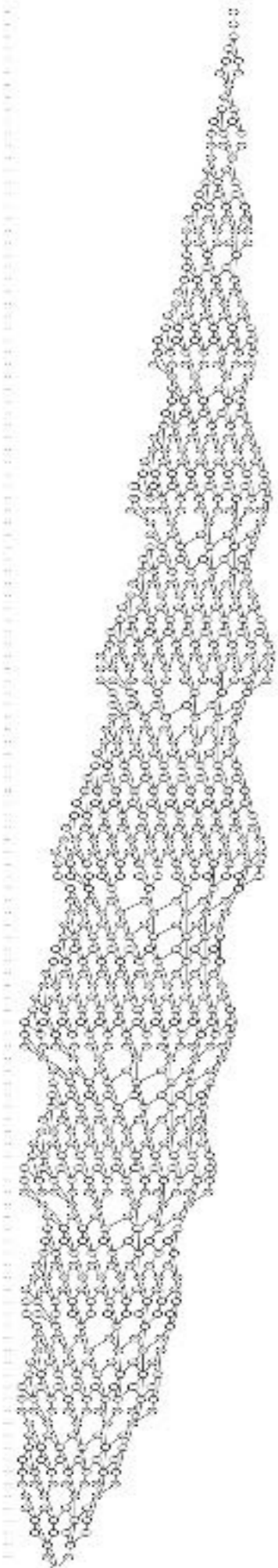
Absorbing dependency

- Two adjacent levels with equal variables can be merged



Factoring algorithm

- Repeat until each p_i , q_j only occurs once in whole BDD:
 - Swap levels repeatedly such that two adjacent levels get equal variables
 - Merge levels
- Any path in resulting BDD gives values of p_i and q_j such that $pq=N$



Only two paths remain

$p=211$ and $q=149$

or

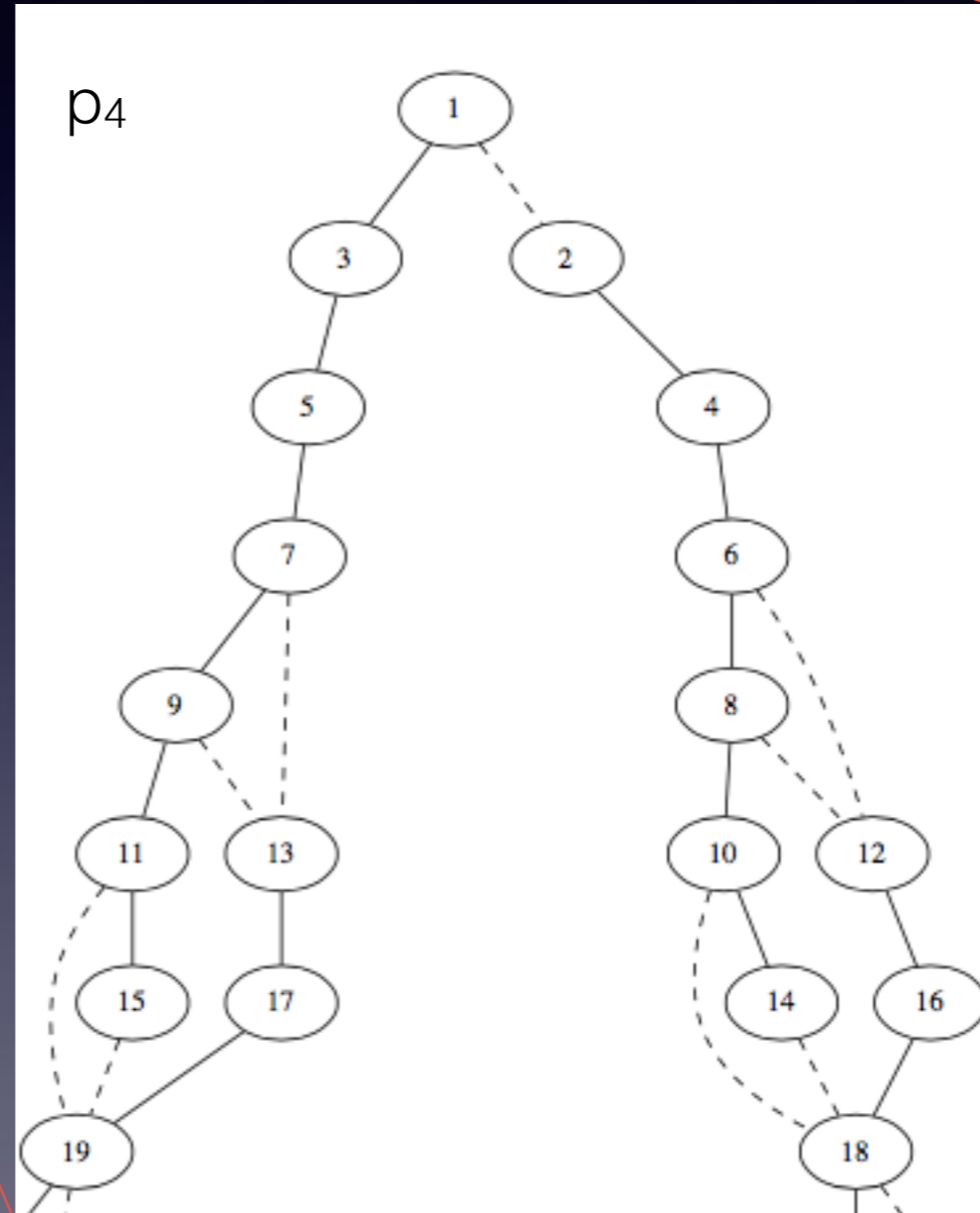
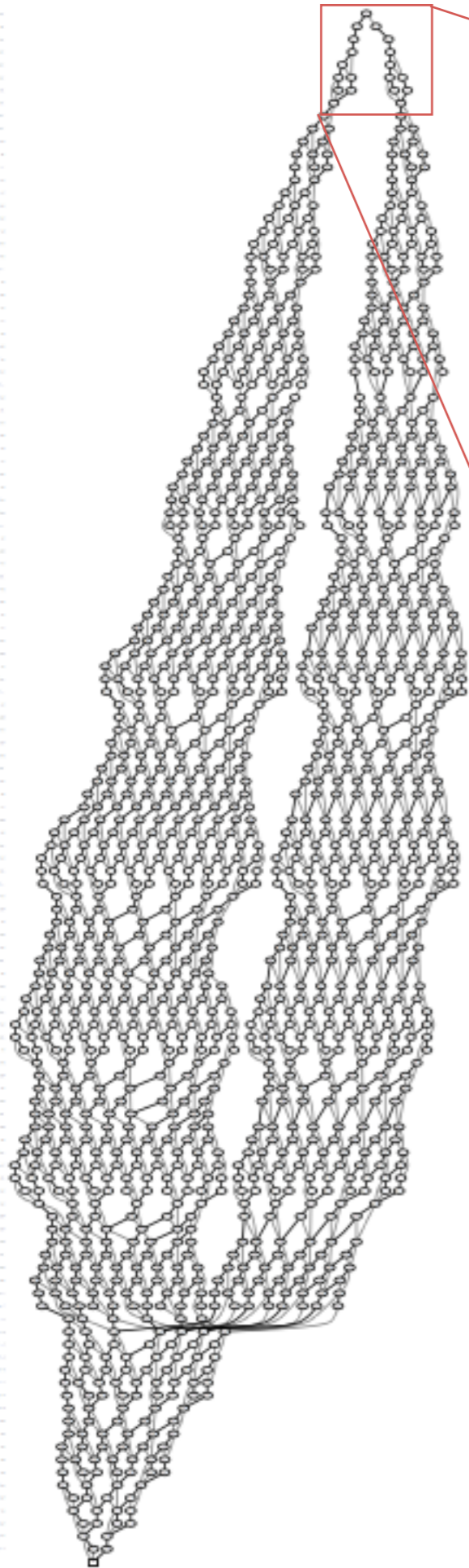
$p=149$ and $q=211$

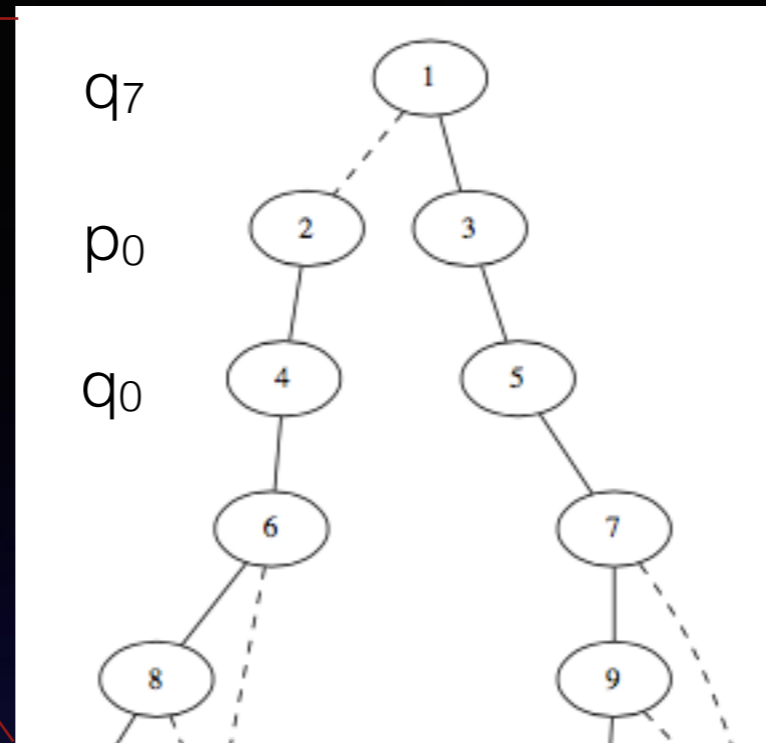
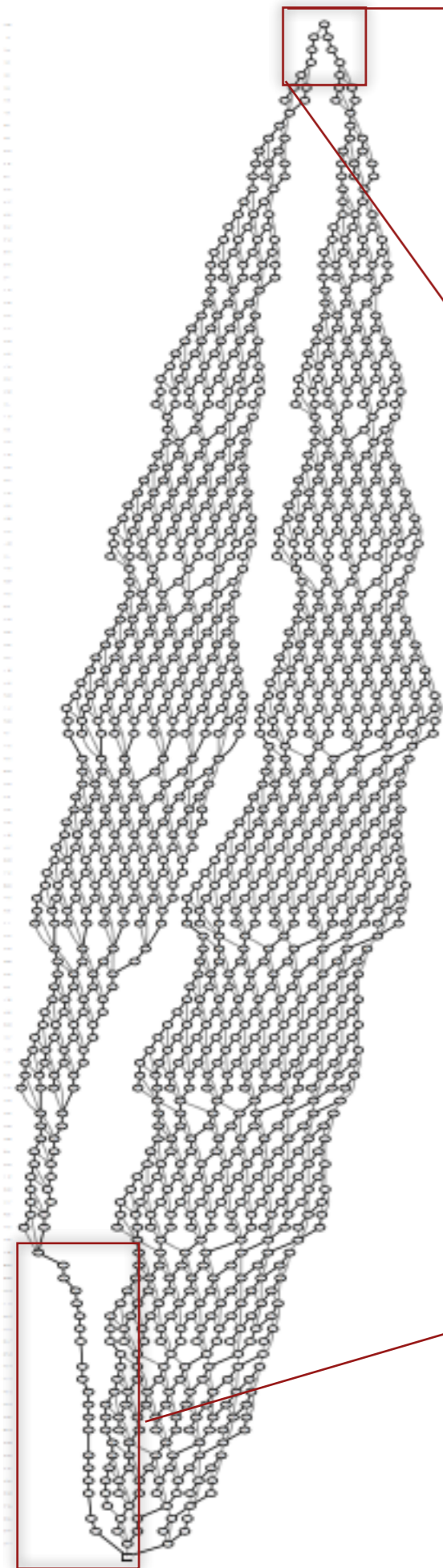
Number of nodes in BDD was never above 891 during solving

Factoring larger numbers?

- Experiments show that simple swap-and-merge algorithm has complexity $\approx \sqrt{N} \approx 2^n$, measured as maximum number of nodes in BDD during solving
- However, BDD representation allows to try other tricks

All levels with p_4 merged and moved to the top





All levels with q_7 merged and moved to the top

$q_7=0$



$q_6=1$ $q_5=1$ $q_4=1$ $q_3=1$
 $p_5=1$ $p_6=1$ $p_7=1$

